

本日のチュートリアルの趣旨

- 今回は、PGAS言語の動向とXcalableMPについての概要だけでなく、Dockerを使ったデモや、post-peta CRESTで研究開発されたGPUクラスタのための拡張 XcalableACC, レンズ実装であるOmni XMPコンパイラの解説、現在、規格部会で議論しているタスク並列への拡張 XMP 2.0まで、以下のスケジュールでご紹介いたします。

本日の予定

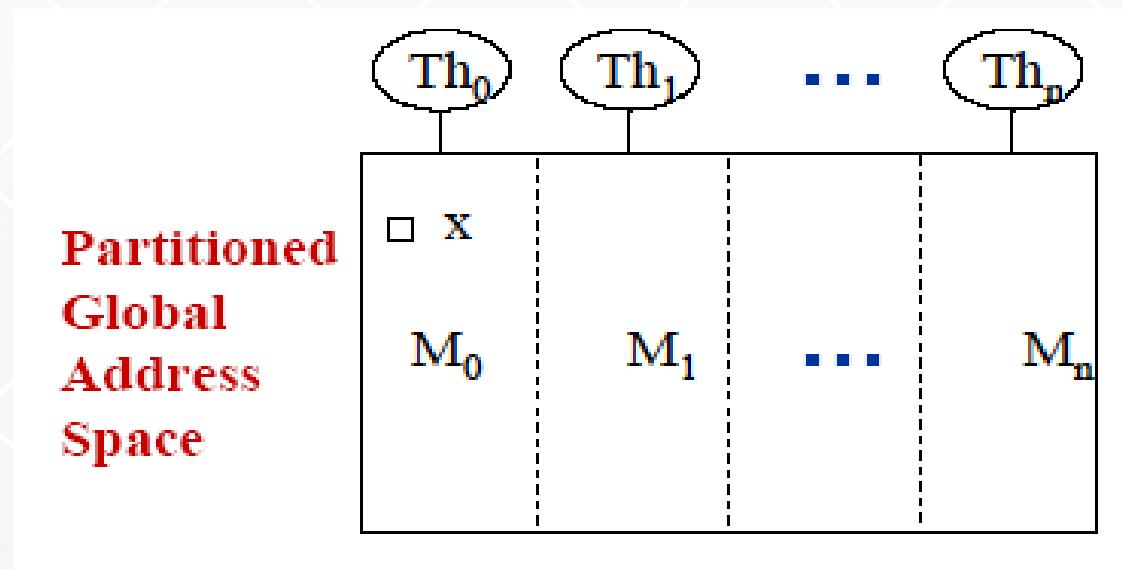
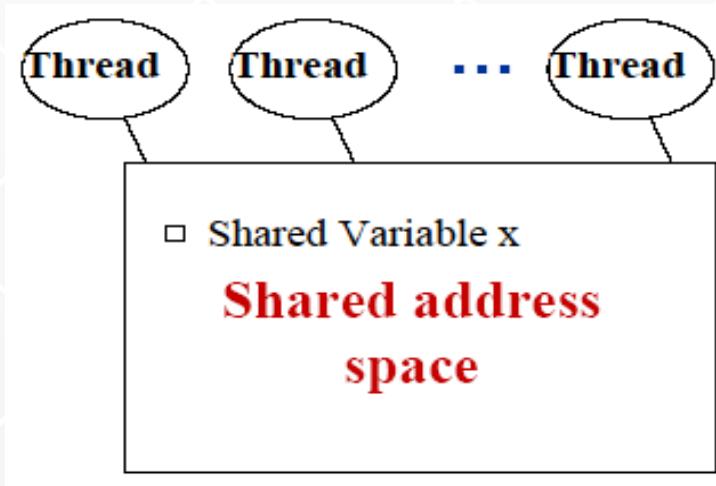
- **13:00~13:20 : PGASの動向とXcalableMP**
 - 佐藤 三久 (理化学研究所)
- **13:20~14:00 : XcalableMPの詳細**
 - 李 珍泌 (理化学研究所)
- **14:00~14:20 : XcalableMPデモ**
 - 中尾 昌広 (理化学研究所)
- **(14:20~14:30 休憩 (コーヒーブレイク))**
- **14:30~15:30 : XcalableACCの概要**
 - 中尾 昌広
- **15:30~16:30 : Omniコンパイラの概要とXMP2.0について**
 - 佐藤 三久

PGASとはなにか？

その動向…

PGAS: Partitioned Global Address Space Programming Model

- ユーザがlocal/globalを宣言する（意識する）
- Partitioned Global Address Space (PGAS) model
- スレッドと分割されたメモリ空間は、対応づけられている（affinity）
 - 分散メモリモデルに対応



Partitioned Global Address Spaceプログラミング言語

- 「shared/global」の発想は、いろいろなところから同時に出てきた。
 - Split-C
 - PC++
 - UPC
 - CAF: Co-Array Fortran
 - X10, Chapel
 - Global Array
 - OpenShmem
- XcalableMP

U P C

- **Unified Parallel C**

- Lawrence Berkeley National Lab.を中心に設計開発

- **Private/Shared を宣言**

- **SPMD**

- MYTHREADが自分のスレッド番号
- 同期機構

- **Barriers**
- **Locks**
- **Memory consistency control**

- **User's view**

- 分割されたshared spaceについて、複数のスレッドが動作する。
- ただし、分割されたshared spaceはスレッドに対してaffinityを持つ。

```
//vect_add.c
#include <upc_relaxed.h>
#define N 100*THREADS
shared int v1[N], v2[N],
v1plusv2[N];

void main(){
    int i;
    for(i=0; i<N; i++)
        if (MYTHREAD==i%THREADS)
            v1plusv2[i]=v1[i]+v2[i];
}
```

```
//vect_add.c
#include <upc_relaxed.h>
#define N 100*THREADS
shared int v1[N], v2[N],
v1plusv2[N];

void main()
{
    int i;
    upc_forall(i=0; i<N; i++;
        v1plusv2[i]=v1[i]+v2[i];
}
```

Affinityが書けるfor構文

UPC: Shared宣言

- Sharedというqualifierを導入する
- Shared array の要素とブロックはthreadのメモリ空間に分散配置される。

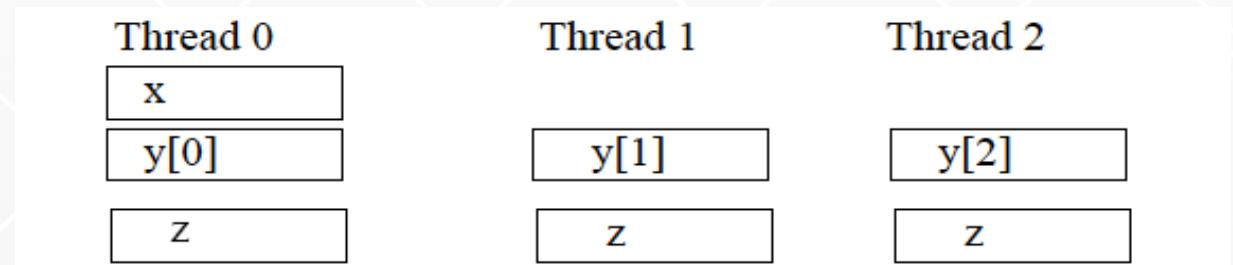
```
shared int x[THREADS] /*One element per thread */  
shared int y[10][THREADS] /*10 elements per thread */
```

- sharedなスカラー変数

```
shared int a;  
/*One item on system (affinity to thread 0) */  
int b; /* one private b at each thread */
```

- Shared Pointer

```
shared int *p;
```

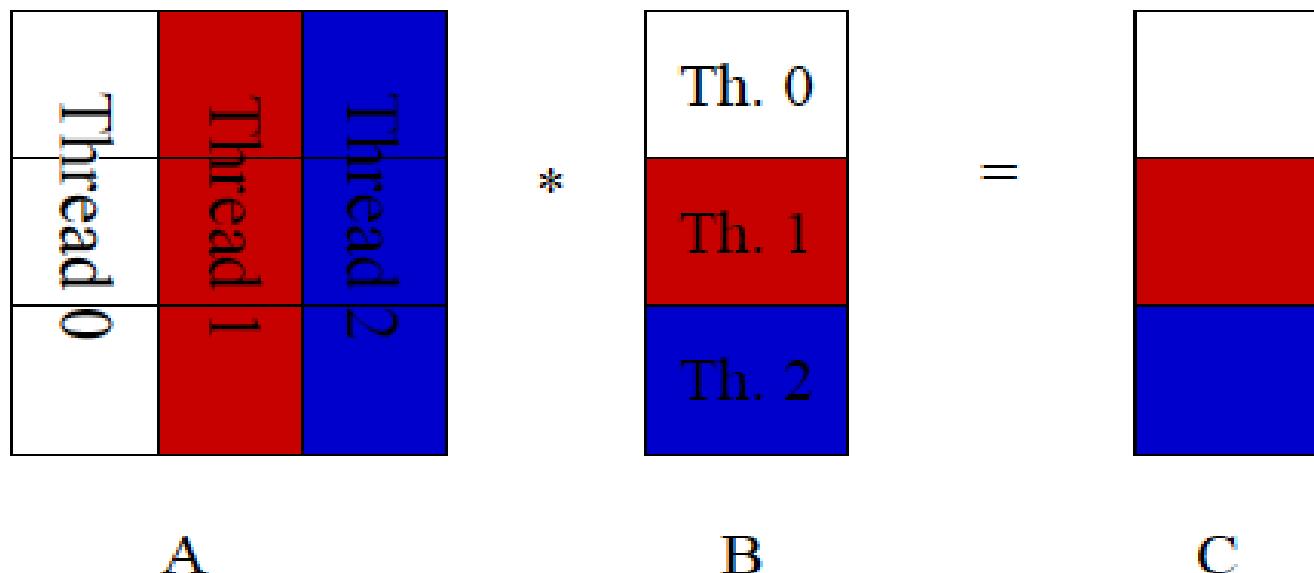


- THREADS = 3の場合

```
shared int x;  
/*x will have affinity to thread 0 */  
shared int y[THREADS];  
int z;
```

行列積の例

```
#include <upc_relaxed.h>
shared int a[THREADS][THREADS];
shared int b[THREADS], c[THREADS];
void main (void) {
    int i, j;
    upc_forall( i = 0 ; i < THREADS ; i++; i) {
        c[i] = 0;
        for ( j= 0 ; j < THREADS ; j++)
            c[i] += a[i][j]*b[j];
    }
}
```



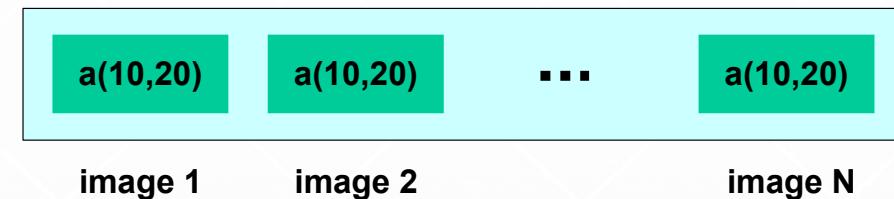
Co-Array Fortran

- Fortran 2008 で正式サポート
- PGAS programming model
 - one-sided communication (GET/PUT)
- SPMD 実行を前提
- Co-array extension
 - 各プロセッサで動くプログラムは、異なる“image”を持つ。

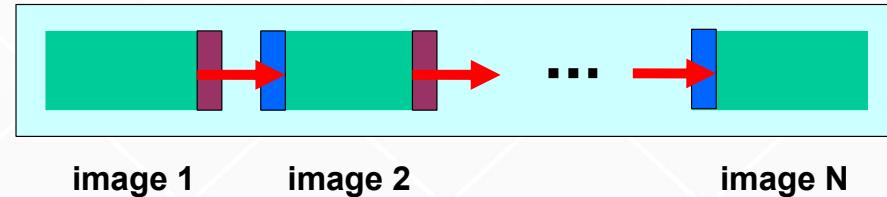
```
real, dimension(n)[:] :: x,y  
x(:) = y(:)[q]
```

qのimageで動くyのデータをローカルなxにコピする(get)
- プログラマは、パフォーマンス影響を与える要素に対して制御する。
 - データの分散配置
 - 計算の分割
 - 通信をする箇所
- データの転送と同期の言語プリミティブをもっている。
 - amenable to compiler-based communication optimization

```
integer a(10,20)[*]
```



```
if (this_image() > 1)  
  a(1:10,1:2) = a(1:10,19:20)[this_image()-1]
```



CAF

- **SPMD process images**
 - fixed number of images during execution
 - images operate asynchronously
- **Both private and shared data**
 - `real x(20, 20)` a private 20x20 array in each image
 - `real y(20, 20)[*]` a shared 20x20 array in each image
- **Simple one-sided shared-memory communication**
 - `x(:,j:j+2) = y(:,p:p+2)[r]` copy columns from image `r` into local columns
- **Synchronization intrinsic functions**
 - `sync_all` – a barrier and a memory fence
 - `sync_mem` – a memory fence
 - `sync_team([team members to notify], [team members to wait for])`
- **Pointers and (perhaps asymmetric) dynamic allocation**

行列積の例

- 2次元のco-Array notationも可能

2次元のCo-dimension

```
real,dimension(n,n)[p,*] :: a,b,c
do k=1,n
    do q=1,p
        c(i,j)[myP,myQ] = c(i,j)[myP,myQ]
            + a(i,k)[myP, q]*b(k,j)[q,myQ]
    enddo
enddo
```

PGAS プログラミングモデルの動向

- ライブラリ的なアプローチが多くなりつつある。
 - Global Array
 - OpenShmem
 - MPI3のput/get (モデルではないが)
 - ...
- UPCも、UPC++でtemplate programming でサポート
- CAFは、正式にFortran 2008でサポート

XcalableMP

- 分散メモリ環境を対象とした指示文ベースの並列言語
- 次世代並列プログラミング言語検討委員会 → 当部会において仕様を検討、提案。
- 2つの並列プログラミングモデルをサポート
 - グローバルビューによる定型的な並列化
 - ローカルビューによる自由度の高い並列化

XcalableMPの現況

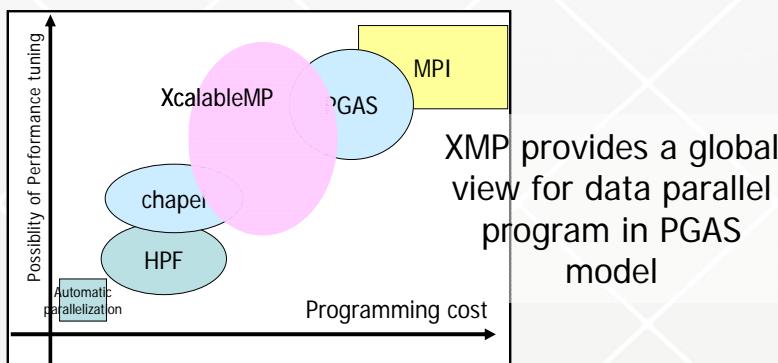
- PCクラスタコンソーシアムで規格を議論
 - 2017/April Version 1.3仕様を公開。
 - Version 1の系統はほぼ「収束」
 - 2015年より、次期仕様「XMP2.0」の検討を開始。
 - PGAS + Multitasking for Multicore
 - Code transformation for Optimization
 - (Accelerator)
- 理研・筑波大で、レファレンス実装
 - Omni XMP コンパイラ

www.xcalablemp.org

omni-compiler.org

XcalableMP(XMP) <http://www.xcalablemp.org>

- What's XcalableMP (XMP for short)?
 - A PGAS programming model and language for distributed memory , proposed by **XMP Spec WG**
 - XMP Spec WG is a special interest group to design and draft the specification of XcalableMP language. It is now organized under **PC Cluster Consortium**, Japan. Mainly active in Japan, but open for everybody.
- Project status (as of June 2016)
 - XMP Spec **Version 1.2.1** is available at XMP site. new features: mixed OpenMP and OpenACC , libraries for collective communications.
 - Reference implementation by U. Tsukuba and Riken AICS: **Version 1.0 (C and Fortran90)** is available for PC clusters, Cray XT and K computer. Source-to-Source compiler to code with the runtime on top of MPI and GasNet.
- **HPCC class 2 Winner 2013. 2014**



- Language Features
 - Directive-based language extensions for Fortran and C for PGAS model
 - Global view programming with global-view distributed data structures for data parallelism
 - SPMD execution model as MPI
 - pragmas for data distribution of global array.
 - Work mapping constructs to map works and iteration with affinity to data explicitly.
 - Rich communication and sync directives such as "gmove" and "shadow".
 - Many concepts are inherited from HPF
- Co-array feature of CAF is adopted as a part of the language spec for local view programming (also defined in C).

```
int array[YMAX][XMAX];  
#pragma xmp nodes p(4)  
#pragma xmp template t(YMAX)  
#pragma xmp distribute t(block) on p  
#pragma xmp align array[i]** to t(i)
```

Code example

```
main(){  
    int i, j, res;  
    res = 0;  
  
    #pragma xmp loop on t(i) reduction(+:res)  
    for(i = 0; i < 10; i++)  
        for(j = 0; j < 10; j++){  
            array[i][j] = func(i, j);  
            res += array[i][j];  
        }  
}
```

data distribution

add to the serial code : incremental parallelization

work sharing and data synchronization

Example of a Global-view XMP Program

- Collaboration in Scale project (with Tomita's Climate Science Team)
- Typical Stencil Code

```
!$xmp nodes p(npz,npz,npz)

!$xmp template (lx,ly,lz) :: t
!$xmp distribute (block,block,block) onto p :: t

!$xmp align (ix,iy,iz) with t(ix,iy,iz) :: 
!$xmp&      sr, se, sm, sp, sn, sl, ...

!$xmp shadow (1,1,1) :: 
!$xmp&      sr, se, sm, sp, sn, sl, ...

...
!$xmp reflect (sr, sm, sp, se, sn, sl)

!$xmp loop (ix,iy,iz) on t(ix,iy,iz)
    do iz = 1, lz-1
    do iy = 1, ly
    do ix = 1, lx
        wu0 = sm(ix,iy,iz ) / sr(ix,iy,iz )
        wu1 = sm(ix,iy,iz+1) / sr(ix,iy,iz+1)
        wv0 = sn(ix,iy,iz ) / sr(ix,iy,iz )
        ...
    
```

declare a node array

declare and distribute a template

align arrays

add shadow area

stencil communication

parallelize loops

Local-view XMP program: Coarray

- XMP includes the coarray feature imported from Fortran 2008 for the local-view programming.
 - Basic idea: data declared as *coarray* can be accessed by remote nodes.
 - Coarray in XMP/Fortran is fully compatible with Fortran 2008.

```
real b(8) [*]  
if (xmp_node_num() == 1) then  
    a(:) = b(:)[2]
```

b is declared as a coarray.

Node 1 gets b from node 2.

- Coarrays can be used in XMP/C.
 - The subarray notation is also available as an extension.

- Declaration
- Put
- Get
- Synchronization

```
float b[8]:[*];  
a[0:3]:[1] = b[3:3]; ← puts b to node 1.  
a[0:3] = b[3:3]:[2]; ← gets b from node 2.  
void xmp_sync_all(int *status)
```

XcalableMP as evolutional approach

- **We focus on migration from existing codes.**
 - Directive-based approach to enable parallelization by adding directives/pragma.
 - Also, should be from MPI code. Coarray may replace MPI.
- **Learn from the past**
 - Global View for data-parallel apps. Japanese community had experience of HPF for Global-view model.
- **Specification designed by community**
 - Spec WG is organized under the PC Cluster Consortium, Japan
- **Design based on PGAS model and Coarray (From CAF)**
 - PGAS is an emerging programming model for exascale!
- **Used as a research vehicle for programming lang/model research.**
 - XMP 2.0 for multitasking.
 - Extension to accelerator (XACC)

XcalableACC(ACC) = XcalableMP+OpenACC+a

● Extension of XcalableMP for GPU

- A part of JST-CREST project “Research and Development on Unified Environment of Accelerated Computing and Interconnection for Post-Petascale” of U. Tsukuba leaded by Prof. Taisuke Boku
- An “orthogonal” integration of XcalableMP and OpenACC
 - Data distribution for both host and GPU by XcalableMP
 - Offloading computations in a set of nodes by OpenACC
- Proposed as unified parallel programming model for many-core architecture & accelerator
 - GPU, Intel Xeon Phi
 - OpenACC supports many architectures

Source Code Example: NPB CG

```
#pragma xmp nodes p(NUM_COLS, NUM_ROWS)
#pragma xmp template t(0:NA-1,0:NA-1)
#pragma xmp distribute t(block, block) onto p
#pragma xmp align w[i] with t(*,i)
#pragma xmp align q[i] with t(i,*)
double a[NZ];
int rowstr[NA+1], colidx[NZ];
...
#pragma acc data copy(p,q,r,w,rowstr[0:NA+1]¥
, a[0:NZ], colidx[0:NZ])
{
    ...
#pragma xmp loop on t(*,j)
#pragma acc parallel loop gang
    for(j=0; j < NA; j++){
        double sum = 0.0;
#pragma acc loop vector reduction(:sum)
        for (k = rowstr[j]; k < rowstr[j+1]; k++)
            sum = sum + a[k]*p[colidx[k]];
        w[j] = sum;
    }
#pragma xmp reduction(:w) on p(:,*) acc
#pragma xmp gmove acc
    q[:] = w[:];
    ...
} //end acc data
```